

# The double category $\mathbb{I}nt(\mathbf{Poly}_+)$ models control flow and data flow

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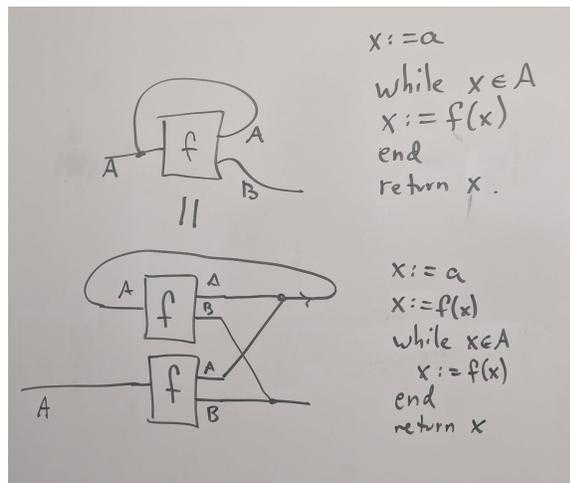
Grigory Kondyrev

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Joint work in progress, with Grigory Kondyrev @ Noeon Research

## I. Introduction

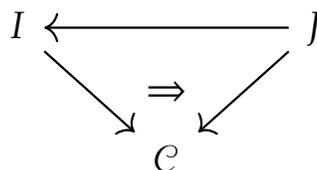
### A. Control flow as while loops (cocartesian traces) 3 mins



### B. Data flow as “query” morphisms

3 mins

1. Morphisms between formal limits within an ontology
2. Roughly  $\text{Diag}(\mathcal{C})$ ,

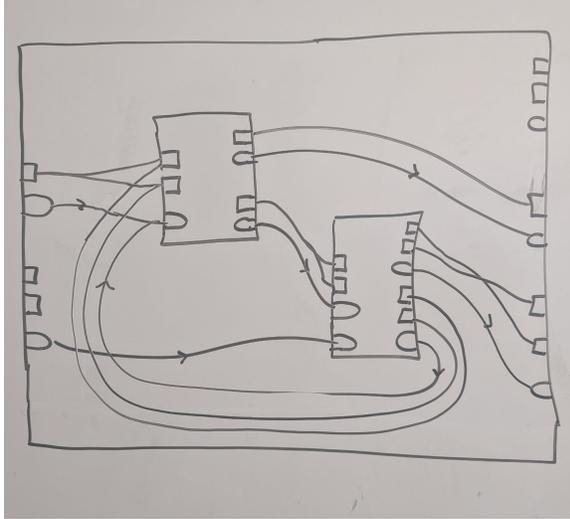


3. e.g., when  $\mathcal{C} = 3$ , we have the query morphism

$$X_1 \times X_2 \times X_3 \rightarrow X_1 \times X_1 \times X_3$$

C. Control and data flow diagrams

3 mins



D. Plan for the talk

1 min

1. Uniform traced monoidal categories  $\mathcal{U}$
2. The thin cocartesian double category  $\mathbf{Int}(\mathcal{U})$
3.  $\mathbf{Int}(\mathbf{Poly}_\star)$  and  $\mathbf{Int}(c\text{-Set}[d]_\star)$ .

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## II. Uniform traced monoidal categories $\mathcal{U}$

A. Traced symmetric monoidal categories

5 mins

1. String diagrams for monoidal categories
2. String diagrams for trace
3. Example:  $(\mathbf{Set}_\star, 0_\star, +_\star)$ ,  $(\mathbf{FDVect}_{\mathbb{R}}, \mathbb{R}, \otimes)$ .
4. Any full (monoidal) subcat'y of a traced cat'y is traced.

B. Uniformity (Hasegawa 2004 *The Uniformity Principle on Traced Monoidal Categories*) 4 mins

$$\begin{array}{ccc}
 A + U & \xrightarrow{f} & B + U \\
 A+h \downarrow & \checkmark & \downarrow B+h \\
 A + V & \xrightarrow{g} & B + V
 \end{array}
 \Rightarrow \text{Tr}_{A,B}^U(f) = \text{Tr}_{A,B}^V(g): A \rightarrow B$$

C. The category  $\mathbf{Set}_\star$  is uniform. 4 mins

D. If  $\mathcal{U}$  is uniform traced, so is  $\mathbf{Fun}(\mathcal{C}, \mathcal{U})$  for any  $\mathcal{C}$ . 8 mins

1. Monoidal structure is pointwise  $(X \otimes Y)(c) := X(c) \otimes Y(c)$ .

2. Trace is also pointwise: given  $f: U \otimes X \Rightarrow U \otimes Y$ ,

$$\mathrm{Tr}_{X,Y}^U(f)_c := \mathrm{Tr}_{X(c),Y(c)}^{U(c)}(f_c).$$

3. So  $\mathbf{Fun}(-, \mathcal{U}): \mathbf{Cat}^{\mathrm{op}} \rightarrow \mathbf{TracedMonCat}$ .

4. Any full subcategory of a uniform traced cat'y is uniform traced.

E. More examples of uniform traced. 5 mins

1.  $(\mathbf{Poly}_\star, 0_\star, +_\star)$  is UT: a full subcategory of  $\mathbf{Fun}(\mathbf{Set}, \mathbf{Set}_\star)$ .

2.  $\mathbf{Set}[\mathcal{D}]_\star \subseteq \mathbf{Fun}(\mathcal{D}\text{-Set}, \mathbf{Set}_\star)$  full subcat'y

3. Example:  $\mathbf{Set}[y_1, \dots, y_L]_\star$ , pointed polynomials in  $L$ -variables, e.g.  $x^2 + \mathbb{N}xy^{\mathbb{R}}z^2 + 1$ .

4.  $\mathcal{C}\text{-Set}[\mathcal{D}]_\star = \mathbf{Fun}(\mathcal{C}, \mathbf{Set}[\mathcal{D}]_\star)$ .

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### III. The thin cocartesian double category $\mathbf{Int}(\mathcal{U})$

A. The  $\mathbf{Int}$  constr'n:  $\mathbf{Int}(\mathcal{J})$  is free CC cat'y on  $\mathcal{J}$ . 6 mins

1. Objects:  $\{(X^i, X^o) \in \mathrm{Ob}(\mathcal{J})^2\}$

2. Morphisms  $\varphi: (X^i, X^o) \rightarrow (Y^i, Y^o)$  in  $\mathbf{Int}(\mathcal{J})$  means

$$Y^i + X^o \rightarrow Y^o + X^i$$

Idea:  $X$  is *inside*  $Y$ . Example,  $X = (0, 0)$  means nothing inside.

3. Composition. Given also  $Z^i + Y^o \rightarrow Z^o + Y^i$ , get  $\varphi \circ \psi$  by

$$\mathrm{Tr}_{Z^i+X^o, Z^o+X^i}^{Y^i} (Z^i+Y^i+X^o \rightarrow Z^i+Y^o+X^i \rightarrow Z^o+Y^i+X^i)$$

B. The double category  $\mathbf{Int}(\mathcal{U})$  for  $\mathcal{U}$  uniform 5 mins

1. Definition of  $\mathbf{Int}(\mathcal{U})$ .
  - a. Tight direction is  $\mathcal{U} \times \mathcal{U}$ .
  - b. Loose arrows are  $\mathbf{Int}(\mathcal{U})$  as above.
  - c. Cells are commuting diagrams

$$\begin{array}{ccc}
 Y_1^i + X_1^o & \xrightarrow{\varphi} & Y_1^o + X_1^i \\
 \downarrow & & \downarrow \\
 Y_2^i + X_2^o & \xrightarrow{\psi} & Y_2^o + X_2^i
 \end{array}$$

2. Thin. They compose horizontally by uniformity.

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#### IV. $\mathbf{Int}(\mathbf{Poly}_\star)$ and $\mathbf{Int}(\mathcal{C}\text{-Set}[\mathcal{D}]_\star)$ .

- A. Consider  $\mathbf{Int}(\mathbf{Poly}_\star)$  and  $\mathbf{Int}(\mathcal{C}\text{-Set}[\mathcal{D}]_\star)$  5 mins

1. Both are cocartesian double categories
2. Compact closed in the horizontal direction.
3. Control and data flow diagrams

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## Bonus material

- B. A factorization system and a weird property 10 mins

1. Vertical maps have an OFS:

$$(X^i, X^o) \xrightarrow{L} (X^i, Y^o) \xrightarrow{R} (Y^i, Y^o).$$

2. When  $\mathcal{U}$  has pushouts and an OFS with extensive right class...
  - a. For example,  $\mathbf{Set}_\star$  has “restriction category” OFS

b.  $X \rightarrow Y$  factors as “domain map”  $X \rightarrow X'$  followed by total function  $X' \rightarrow Y$

3. ...we get the following property. Any diagram

$$\begin{array}{ccc}
 (X_1^i, X_1^o) & \twoheadrightarrow & (Y_1^i, Y_1^o) \\
 L \downarrow & & \downarrow R \\
 (X_1^i, X_2^o) & \Downarrow & (Y_2^i, Y_1^o) \\
 \downarrow & & \downarrow \\
 (X^i, X^o) & \twoheadrightarrow & (Y^i, Y^o)
 \end{array}$$

can be factored as

$$\begin{array}{ccc}
 (X_1^i, X_1^o) & \twoheadrightarrow & (Y_1^i, Y_1^o) \\
 L \downarrow & & \downarrow R \\
 (X_1^i, X_2^o) & \Downarrow & (Y_2^i, Y_1^o) \\
 R \downarrow & & \downarrow L \\
 (X_2^i, X_2^o) & \twoheadrightarrow & (Y_2^i, Y_2^o) \\
 \downarrow & \Downarrow & \downarrow \\
 (X^i, X^o) & \twoheadrightarrow & (Y^i, Y^o)
 \end{array}$$

4. Indeed, the top diagram boils down to

$$\begin{array}{ccc}
 Y_1^i + X_1^o & \longrightarrow & Y_1^o + X_1^i \\
 \downarrow & & \downarrow \\
 Y_2^i + X_2^o & & \\
 \downarrow & & \downarrow \\
 Y^i + X^o & \longrightarrow & Y^o + X^i
 \end{array}$$

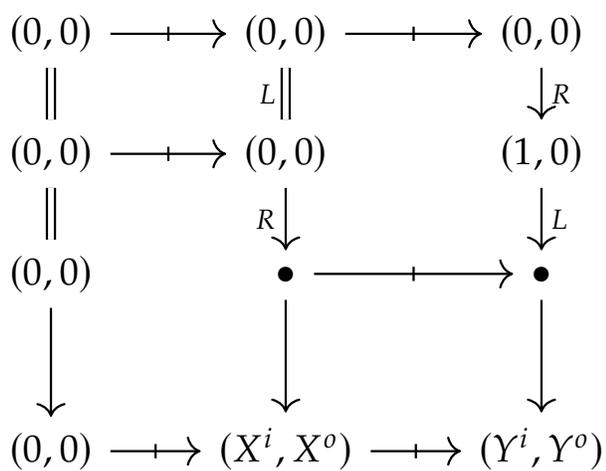
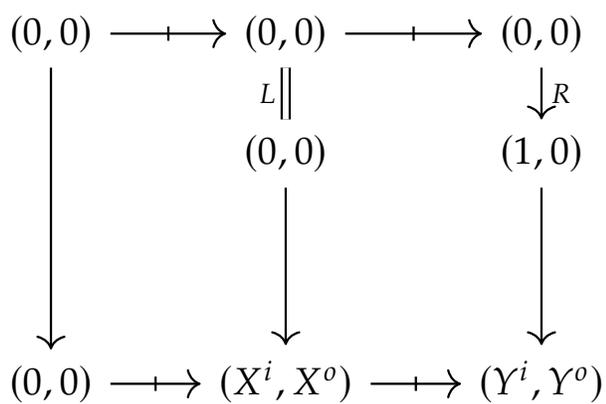
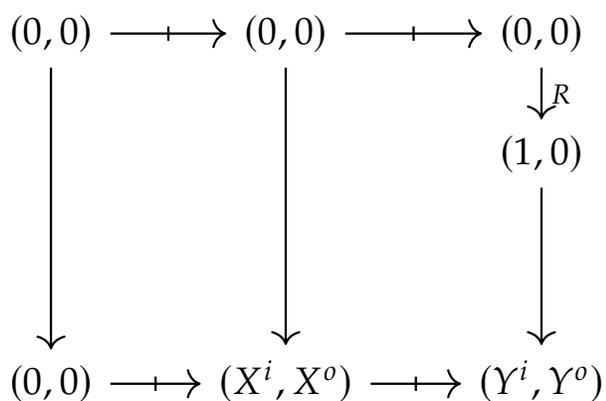
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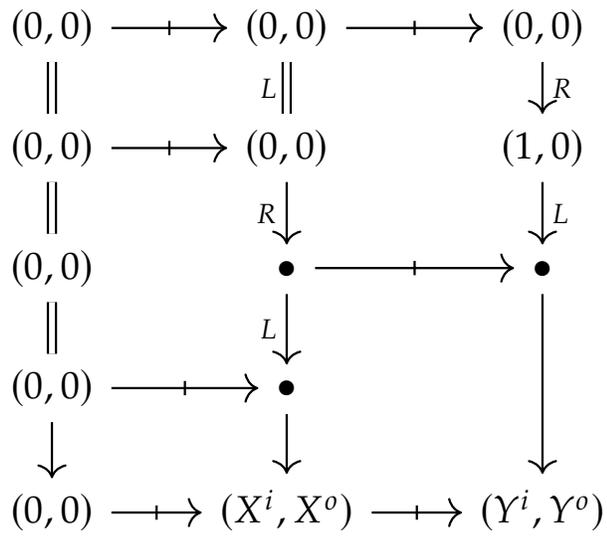
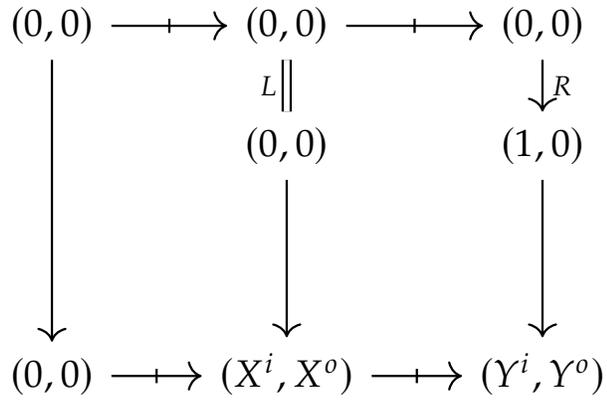
- a. Perform a pushout
- b. Factor the right map

c. Use extensivity to form new pair of objects.

C. Use the above to get trajectories

4 mins





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